Resource-Guided Program Synthesis

Tristan Knoth[1], Di Wang[2], Nadia Polikarpova[1], Jan Hoffmann[2]

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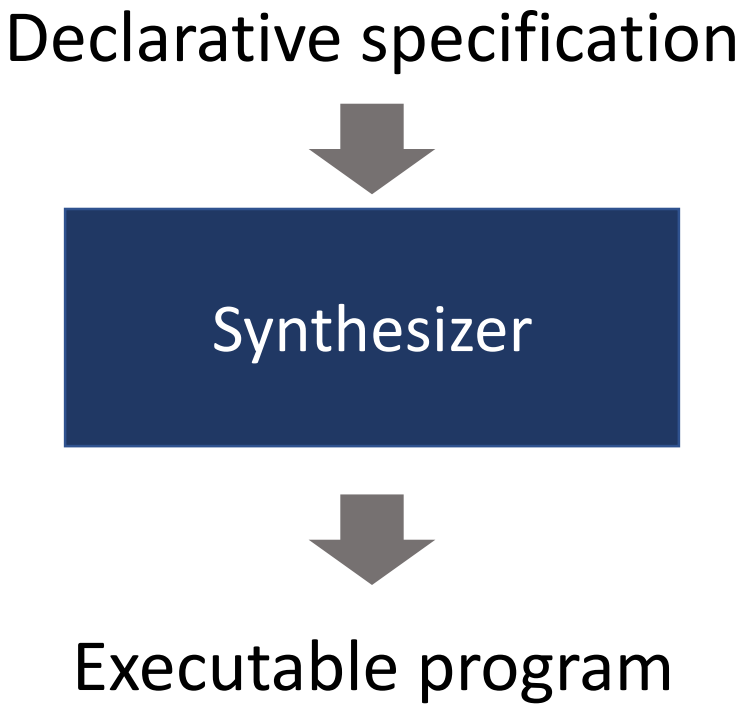
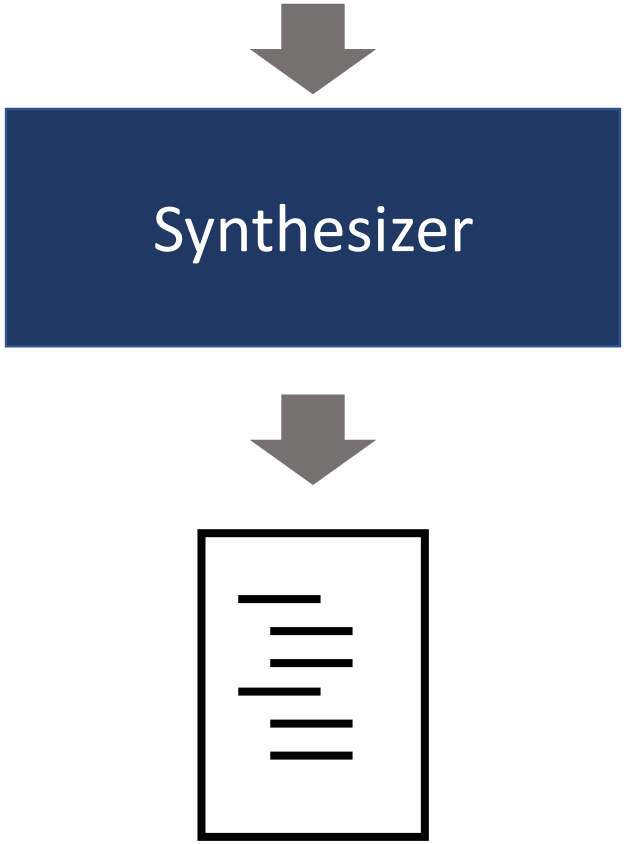
PLDI 2019

Program Synthesis



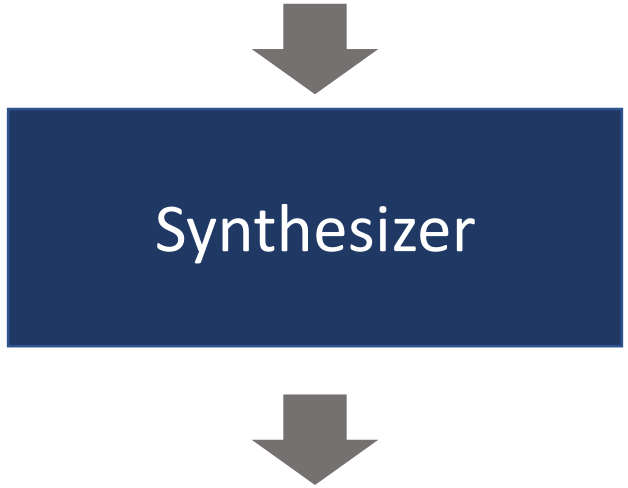
State of the art

“Find the intersection
of two **sorted** lists”



Type-directed synthesis

```
common :: xs: SList a → ys: SList a
        → v: {List a | elems v = elems xs ∩ elems ys}
```




```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(member x ys)
        then common xt ys
        else Cons x (common xt ys)
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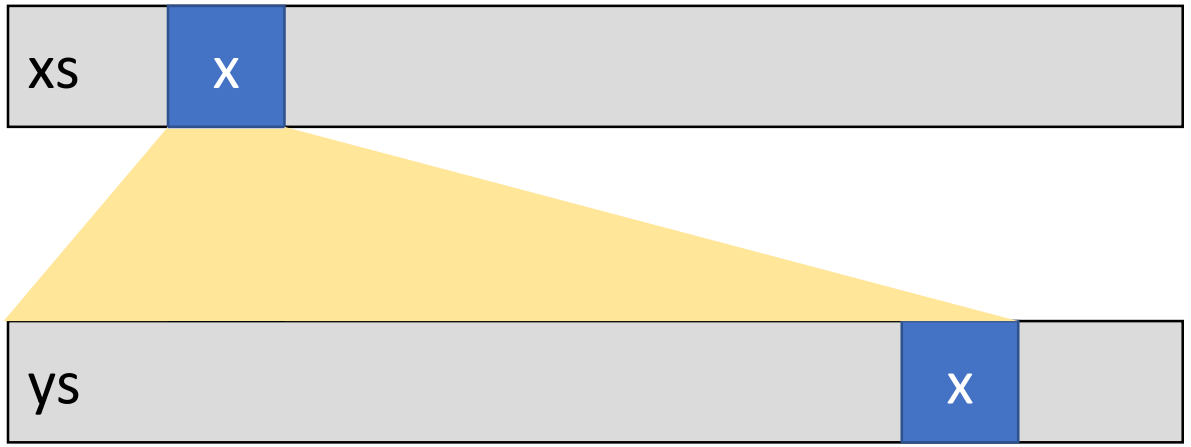

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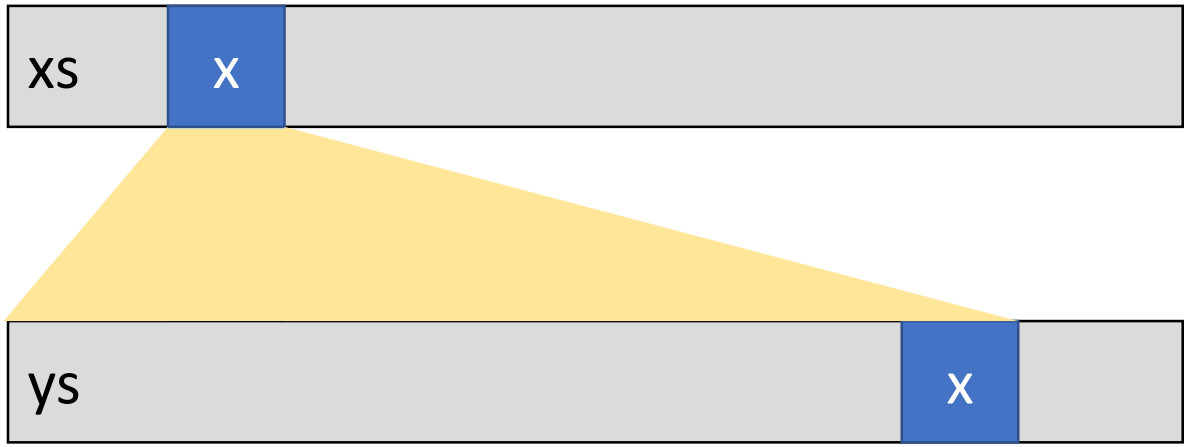

```
common = λ xs. λ ys.
  match xs with
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          if x < y
            then common xt ys
            else if y < x
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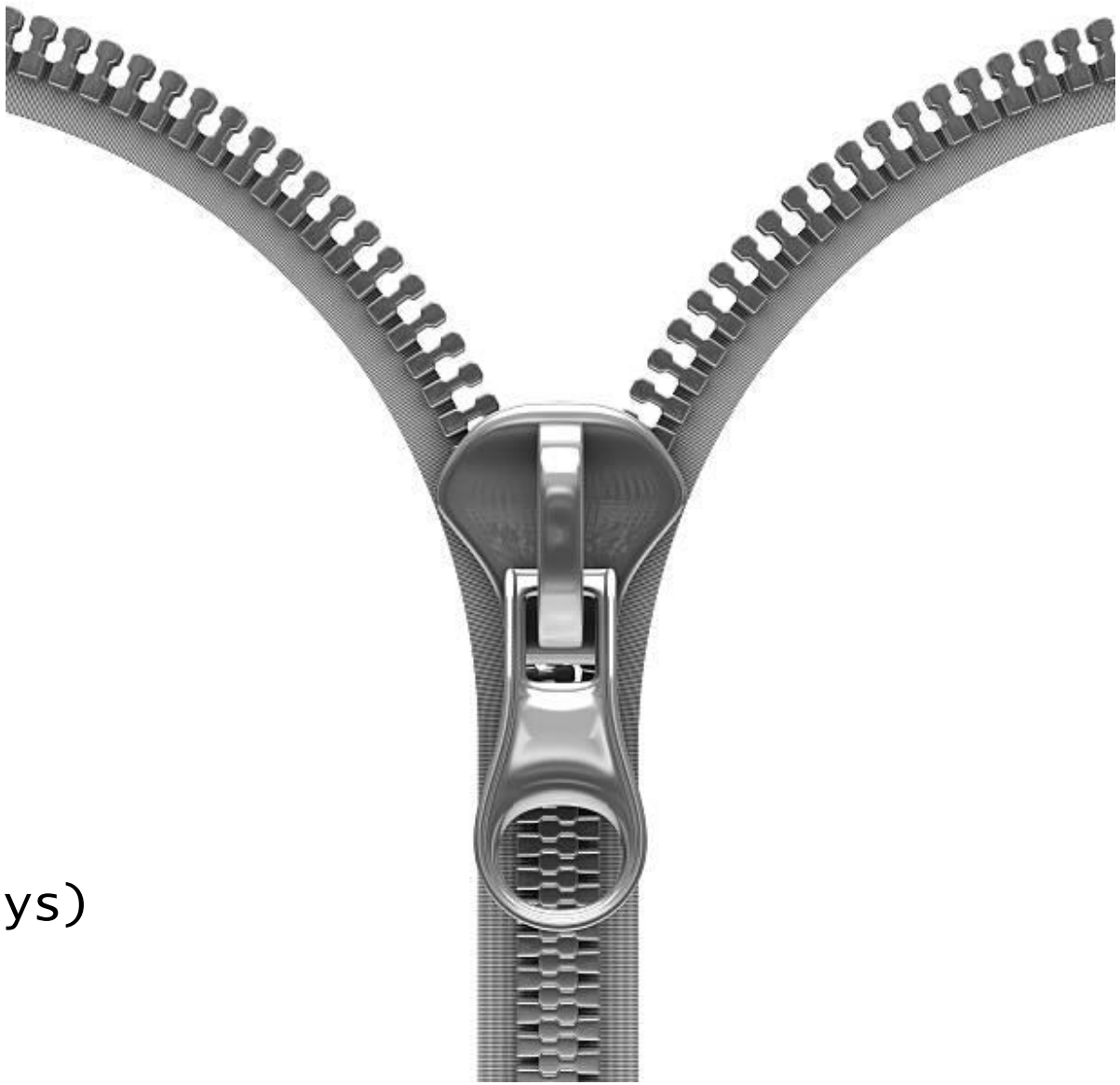
```



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```

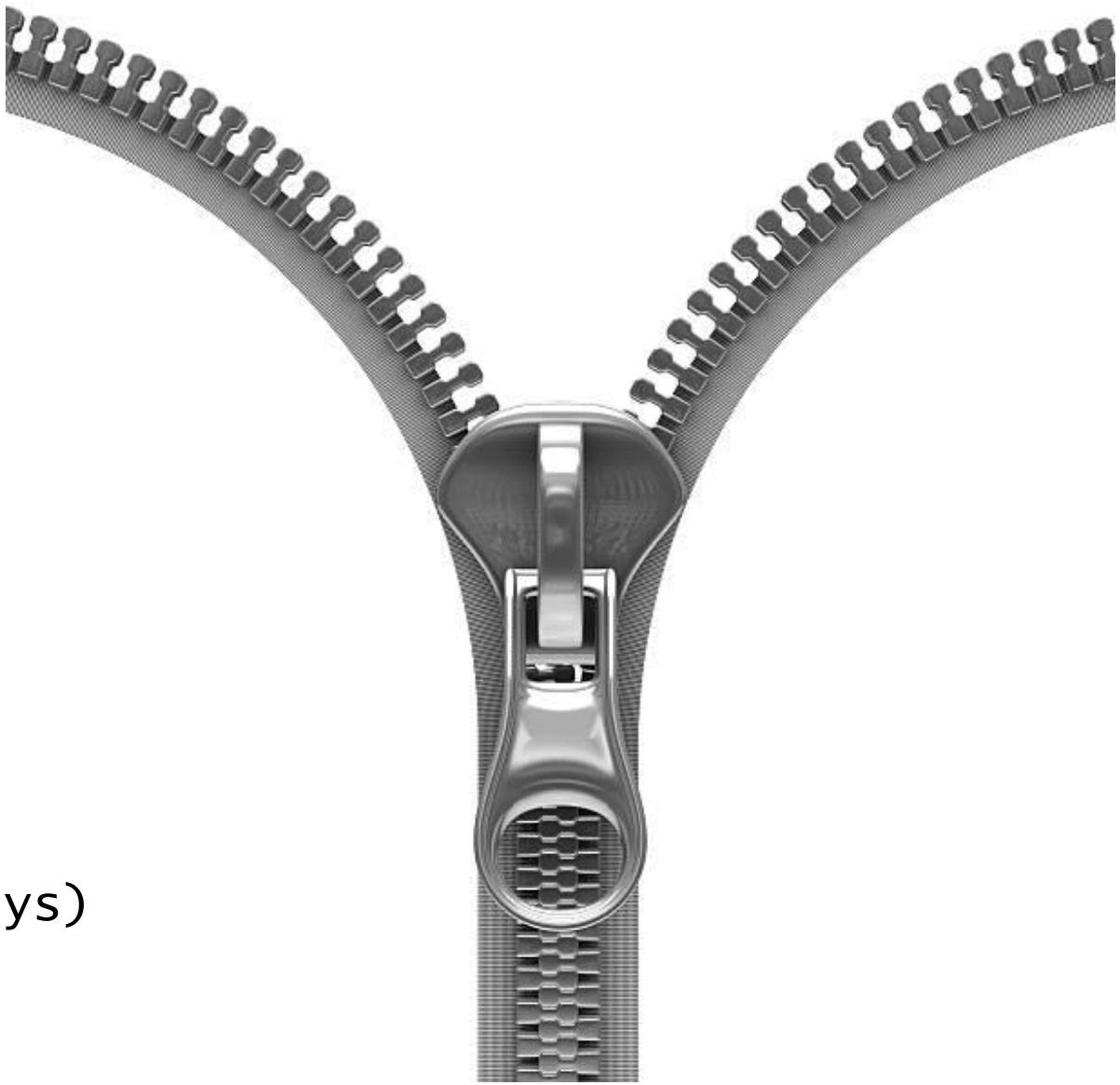

O(m·n)

```
common = λ xs. λ ys.
  match xs with
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```

O(m + n)

```
common = λ xs. λ ys.
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        Nil → Nil
        Cons y yt →
          if x < y
            then common xt ys
            else if y < x
              then common xs yt
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```
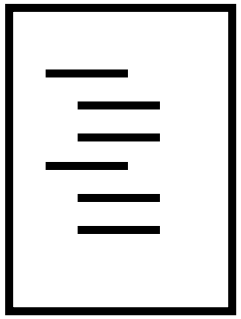

What we have

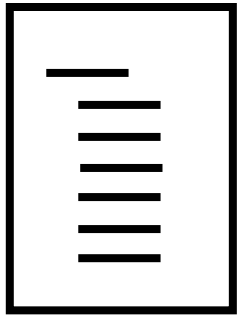
"Find the intersection of two sorted lists"



O(m·n)

What we want

"Find the intersection of two sorted lists **in linear time**"



O(m+n)

ReSyn

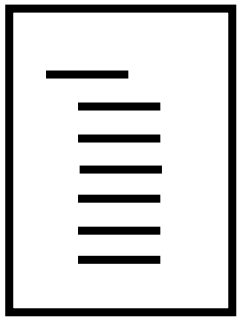
The first resource-aware synthesizer for recursive programs

This talk

1. Specification

"Find the intersection of two sorted lists in linear time"

Synthesizer

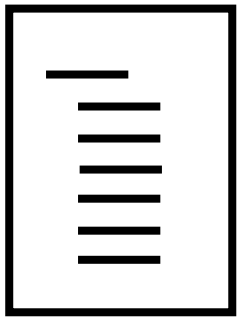


This talk

1. Specification

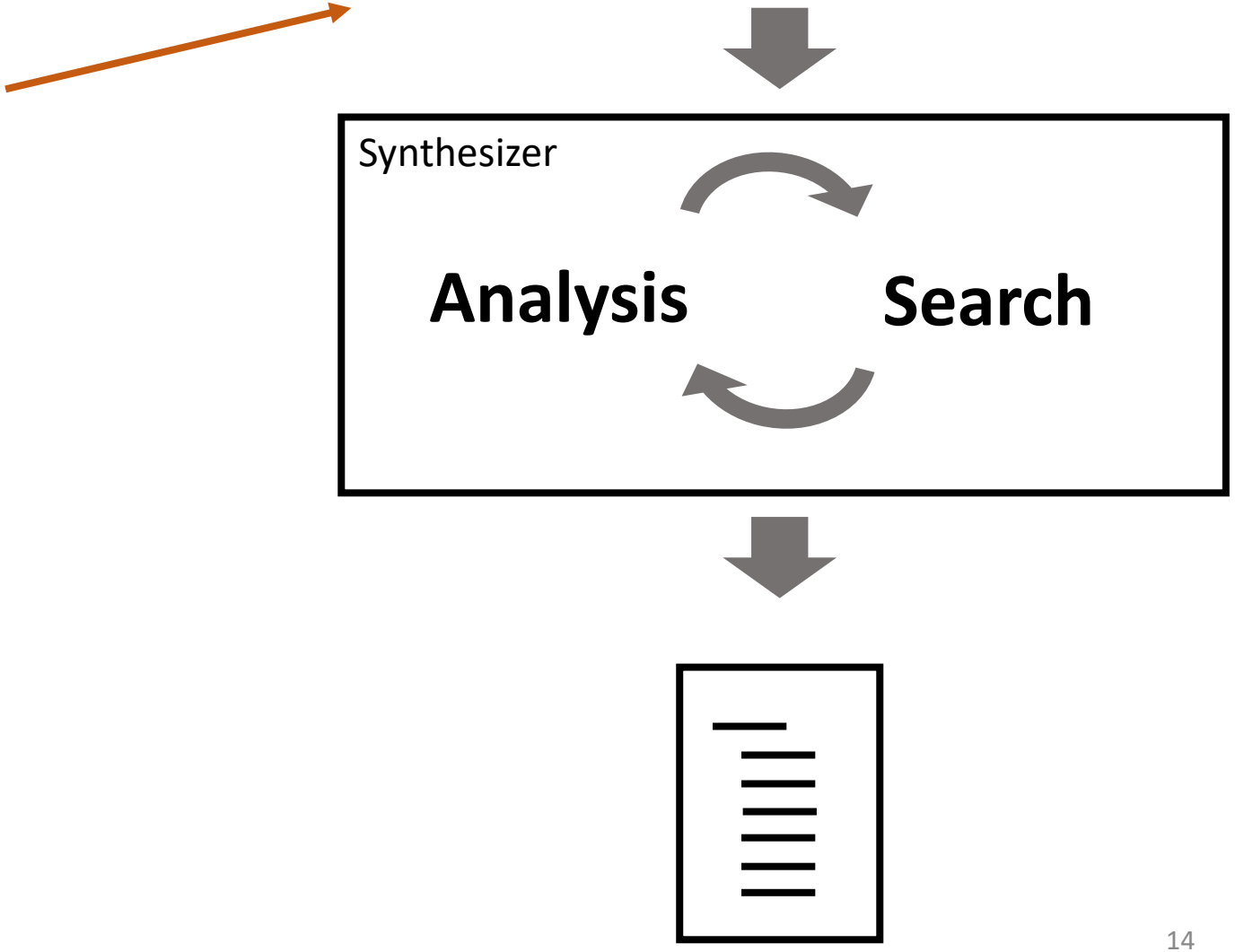
"Find the intersection of two sorted lists **in linear time**"

Synthesizer



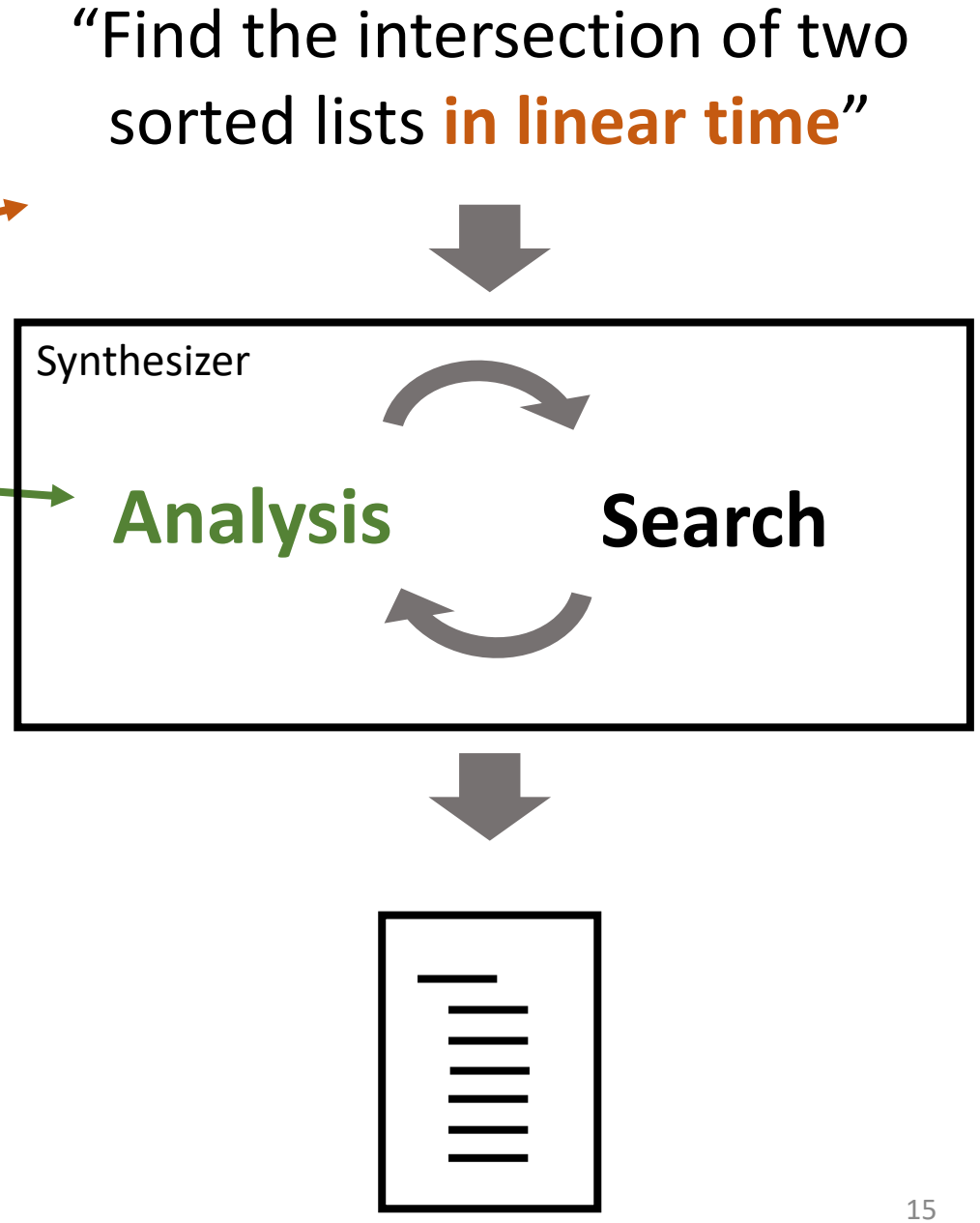
This talk

1. Specification



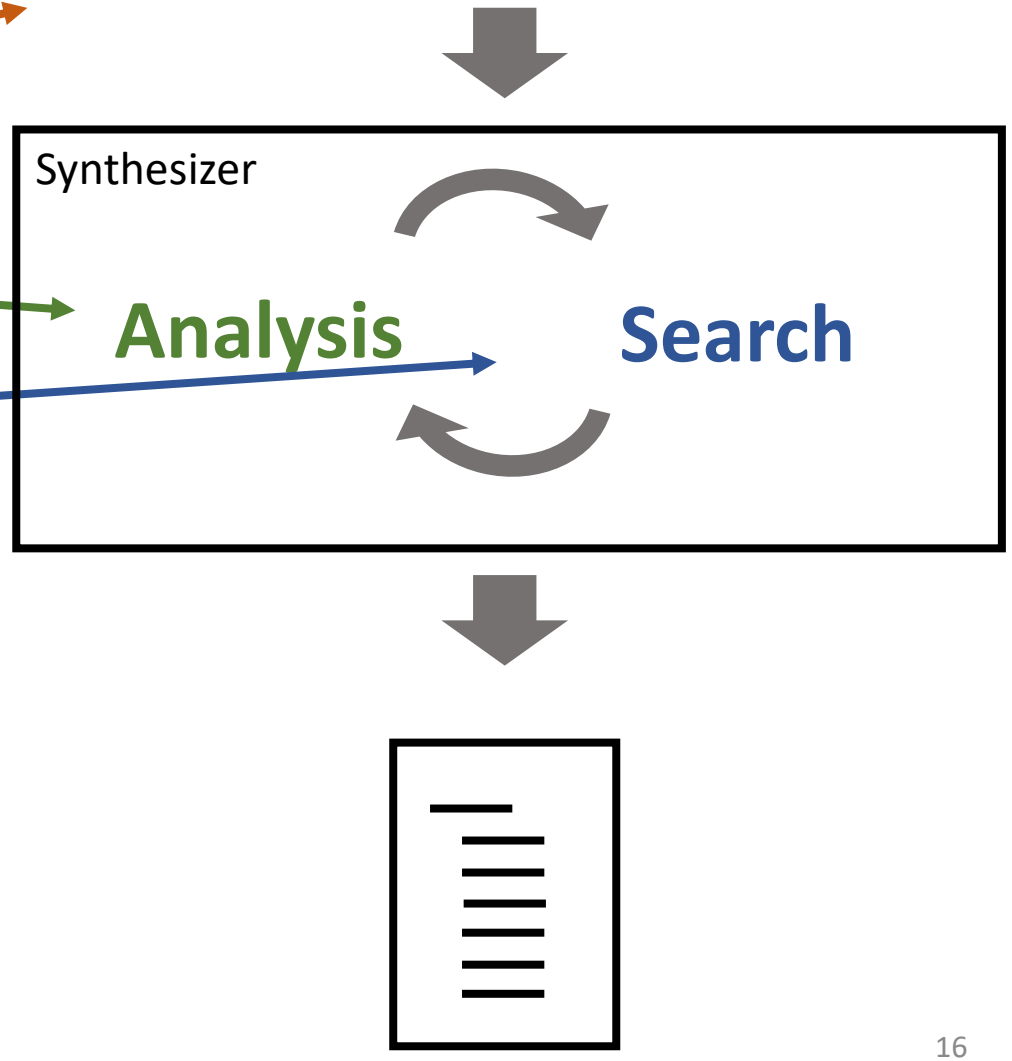
This talk

1. Specification
2. Analysis



This talk

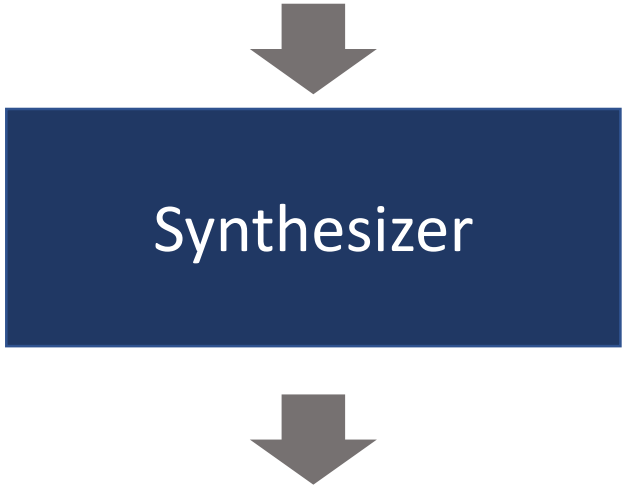
1. Specification
2. Analysis
3. Search



This talk

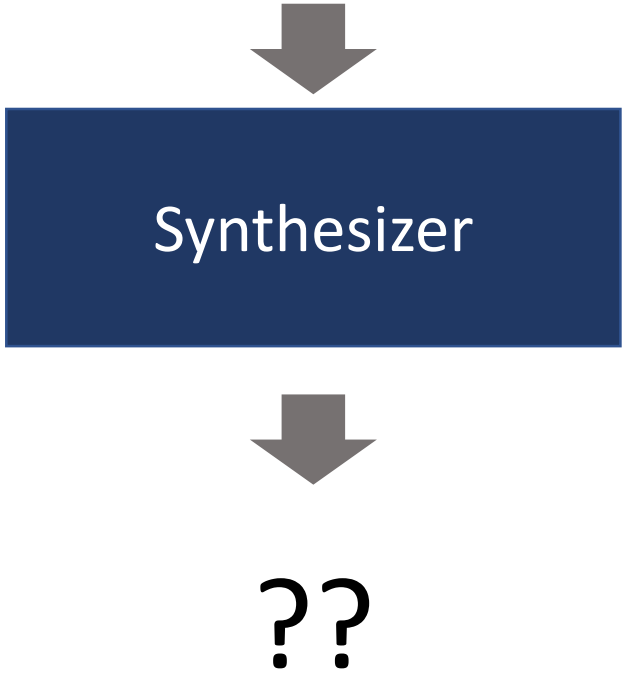
1. Specification
2. Analysis
3. Search

“Find the intersection of two sorted lists in linear time”



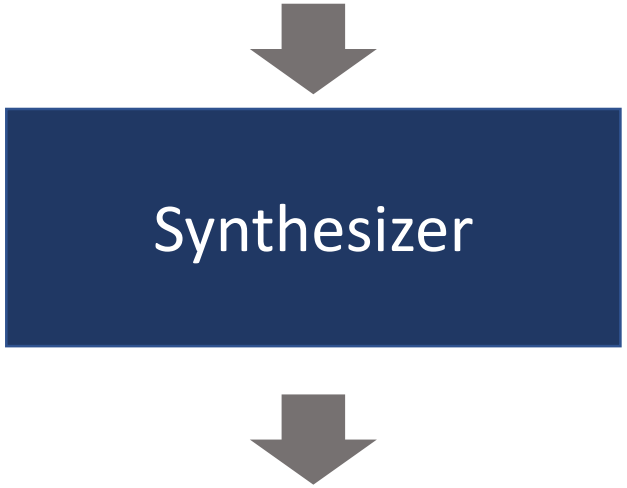
??

“**Find the intersection of two sorted lists** in linear time”



Refinement types

“**Find the intersection of two sorted lists in linear time**”



??

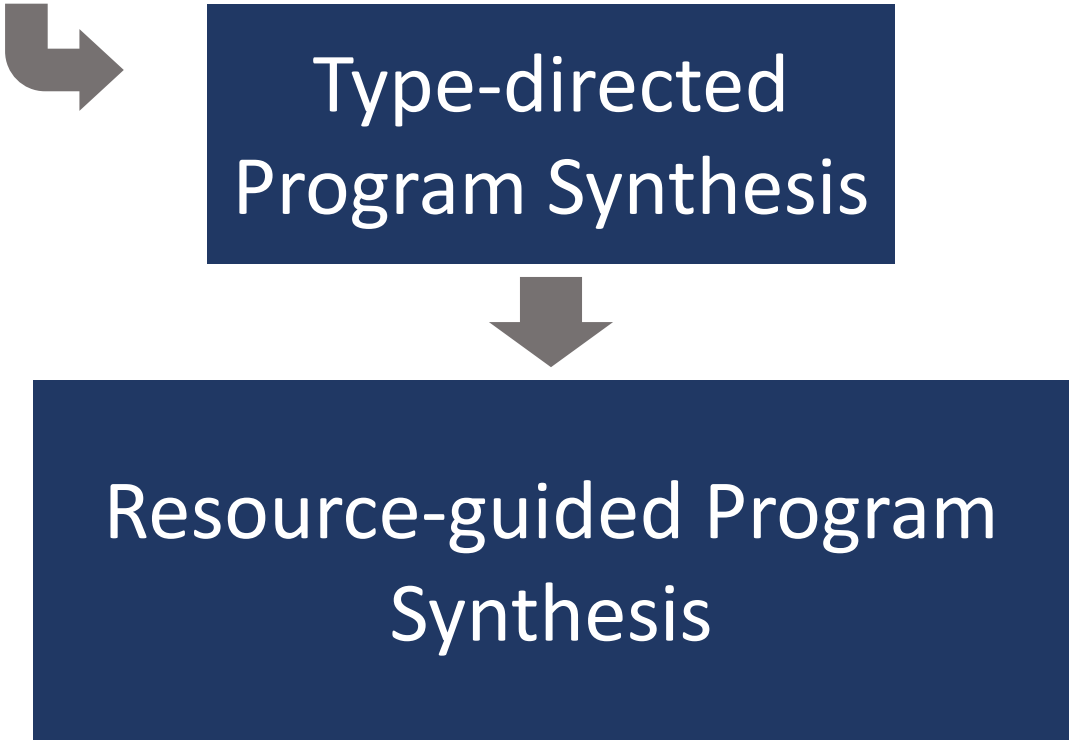
Refinement types

with

Resource annotations

Refinements:

Synquid



[Polikarpova et. al 2016]

Refinements:
Synquid

Resource annotations:
Automated Amortized Resource Analysis

[Polikarpova et. al 2016]

[Hoffmann et al. 2010]

“**Find the intersection of two sorted lists** in linear time”

$$v : \{\mathtt{Int} \mid v \geq 0\}$$

Refinement types

```
common = ??
```

Refinement types

```
common :: xs: SList a → ys: SList a
        → ν: {List a | elems ν = elems xs ∩ elems ys}
common = ??
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Refinement types

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Refinement types

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common :: xs: SList a → ys: SList a
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```
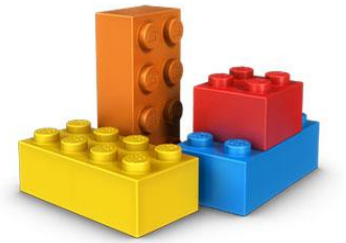

Functional specification

Synquid

Library functions



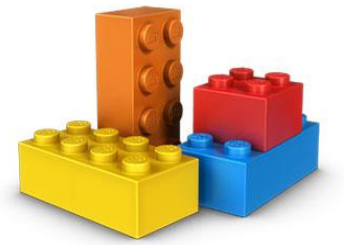
[Polikarpova et. al, 2016]

Functional specification

Synquid

```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(member x ys)
        then common xt ys
        else Cons x (common xt ys)
```

Library functions

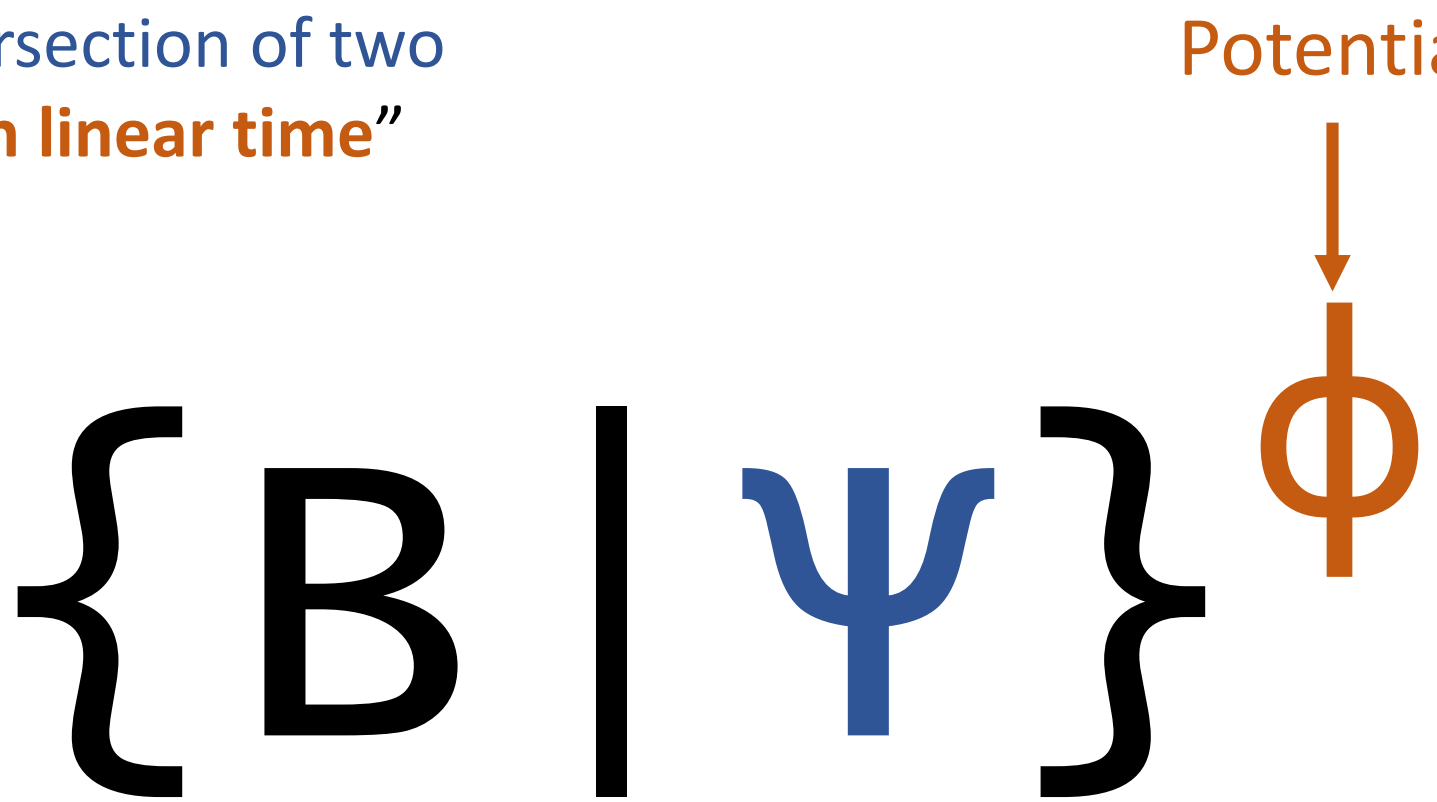


[Polikarpova et. al, 2016]

“Find the intersection of two sorted lists **in linear time**”

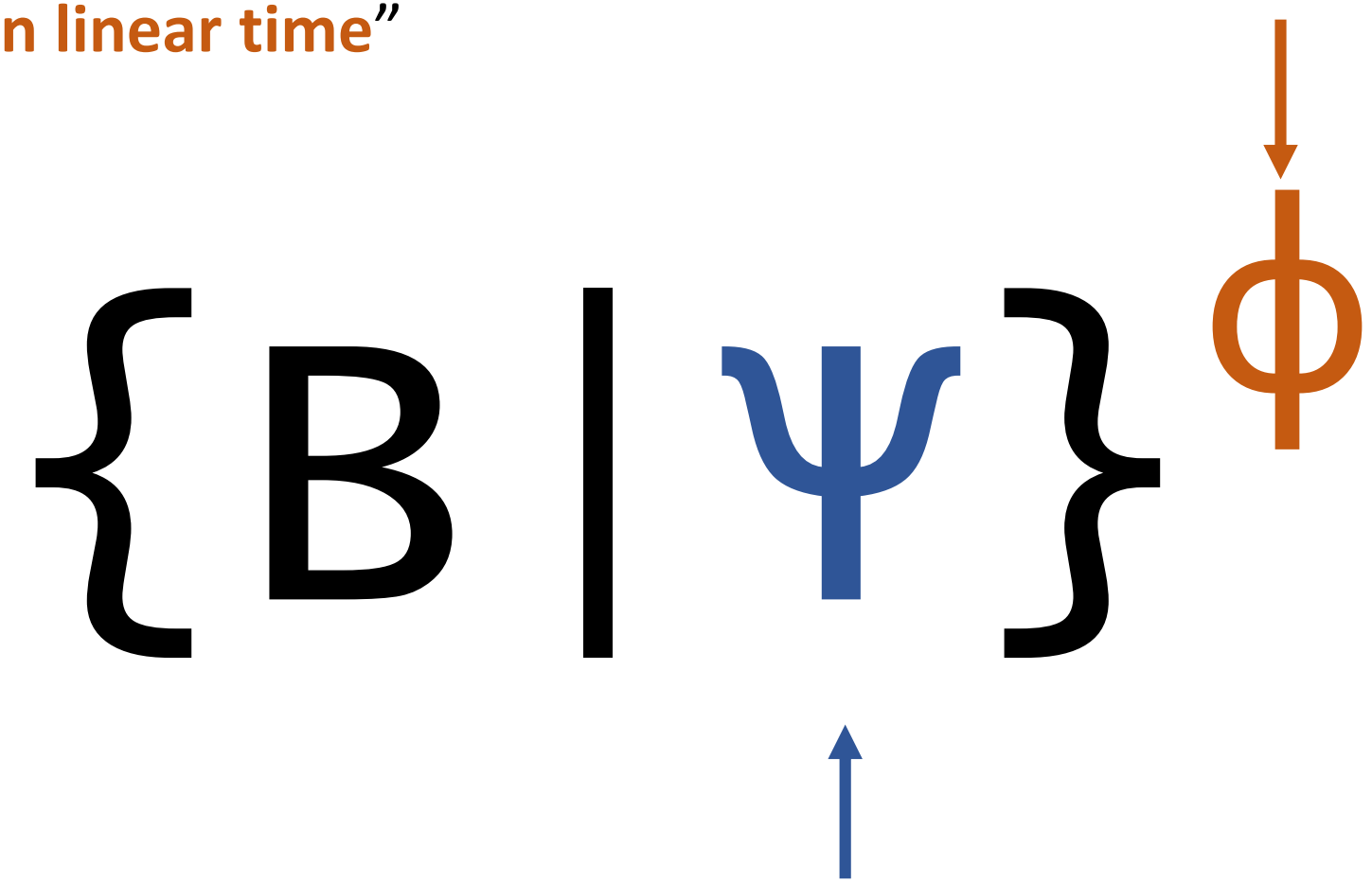
"Find the intersection of two sorted lists **in linear time**"

Potential



"Find the intersection of two sorted lists **in linear time**"

Potential: numeric



Resource annotations

```
common :: xs: SList a → ys: SList a
        → ν: {List a | elems ν = elems xs ∩ elems ys}
common = ??
```

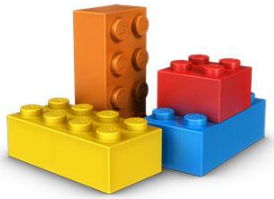
Resource budget

```
common :: xs: SList a¹ → ys: SList a¹
        → ν: {List a | elems ν = elems xs ∩ elems ys}
common = ??
```

Synthesize with ReSyn

```
common :: xs: SList a¹ → ys: SList a¹
        → ν: {List a | elems ν = elems xs ∩ elems ys}
common = ??
```

```
member
Cons, Nil, …
≤, =, !, …
```



Components: member

```
member :: z:a → zs: List a
        → ν:{Bool|ν = (x ∈ elems xs)}
```

Components: member

```
member :: z:a → zs: List a¹
        → ν:{Bool|ν = (x ∈ elems xs)}
```

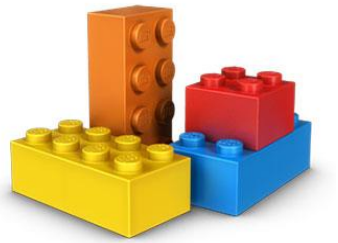
Components: `member`

```
member :: z:a → zs: List a[1]
        → ν:{Bool|ν = (x ∈ elems xs)}
```

Functional specification

Resource bound

Library functions

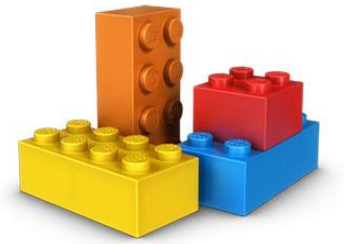


ReSyn

Functional specification

Resource bound

Library functions



ReSyn

```
common = λ xs. λ ys.
  match xs with
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    Cons x xt →
      match ys with
        Nil → Nil
        Cons y yt →
          if x < y
            then common xt ys
            else if y < x
              then common xs yt
              else Cons x (common xs ys)
```

This talk

1. Specification
2. Analysis
3. Search

How do we know **common** does not run in linear time?

```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(member x ys)
        then common xt ys
        else Cons x (common xt ys)
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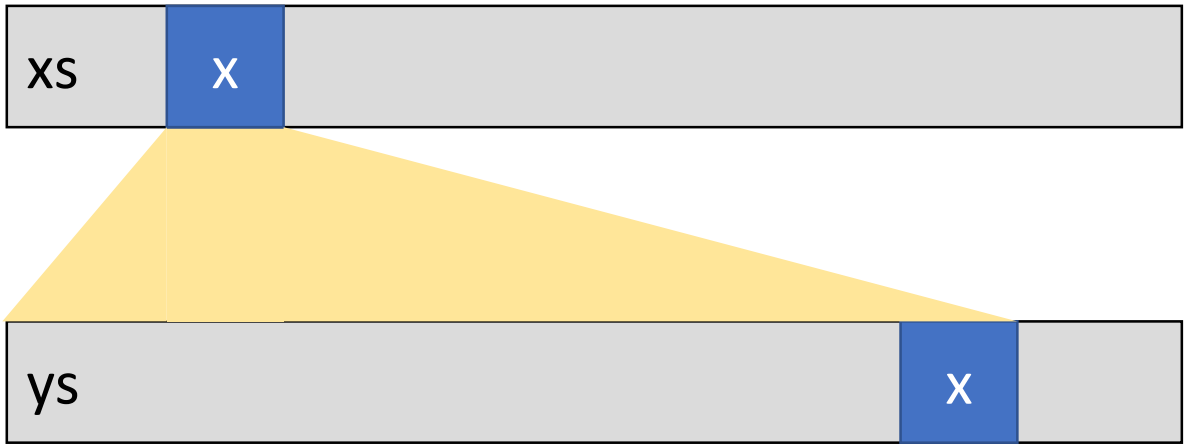


```
member :: z:a → zs: List a¹
        → v:{Bool|v = (x ∈ elems xs)}
```

How do we automate this reasoning?

```
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```

```
common :: xs: SList a¹ → ys: SList a¹ → v: {List a |…}
common = λ xs. λ ys.
  match xs with
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      if !(member x ys)
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```

Can we partition the allotted resources between all function calls?

```
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```

ys :: SList a^1

```
common = λ xs. λ ys.
  match xs with
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```

```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(member x (ys :: List aᵖ))
        then common xt (ys :: List aᵠ)
        else Cons x (common xt ys)
```

```
common = λ xs. λ ys.
  match xs with
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```

List a^p <: List a^1

$$\frac{a <: b \qquad p \geq q}{a^p <: b^q}$$

$$\frac{}{\texttt{List}\ a^p <: \texttt{List}\ b^q}$$

```
member :: z:a → zs: List a¹ → ν:{Bool|…}
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```

$$\text{List } a^{p} <: \text{List } a^{1}$$
$$p \geq 1$$

```
common :: xs: SList a¹ → ys: SList a¹ → v: {List a |…}
```

```
common = λ xs. λ ys.
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    Cons x xt →
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```
common :: xs: SList a¹ → ys: SList a¹ → v: {List a |…}
```

$$\text{List } a^q <: \text{List } a^1 \qquad q \geq 1$$

```
common = λ xs. λ ys.
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        then common xt (ys :: List aᵠ)
        else Cons x (common xt ys)
```

Sharing ⟶ SList a^1 ⟶ SList a^p, SList a^q

```
common = λ xs. λ ys.
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```

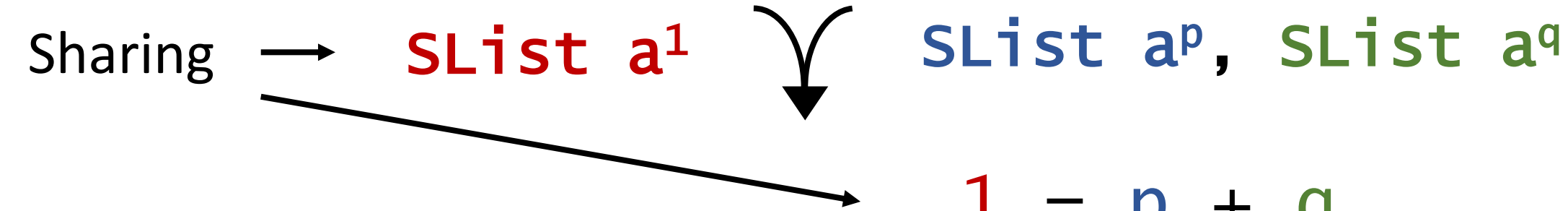


```
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```
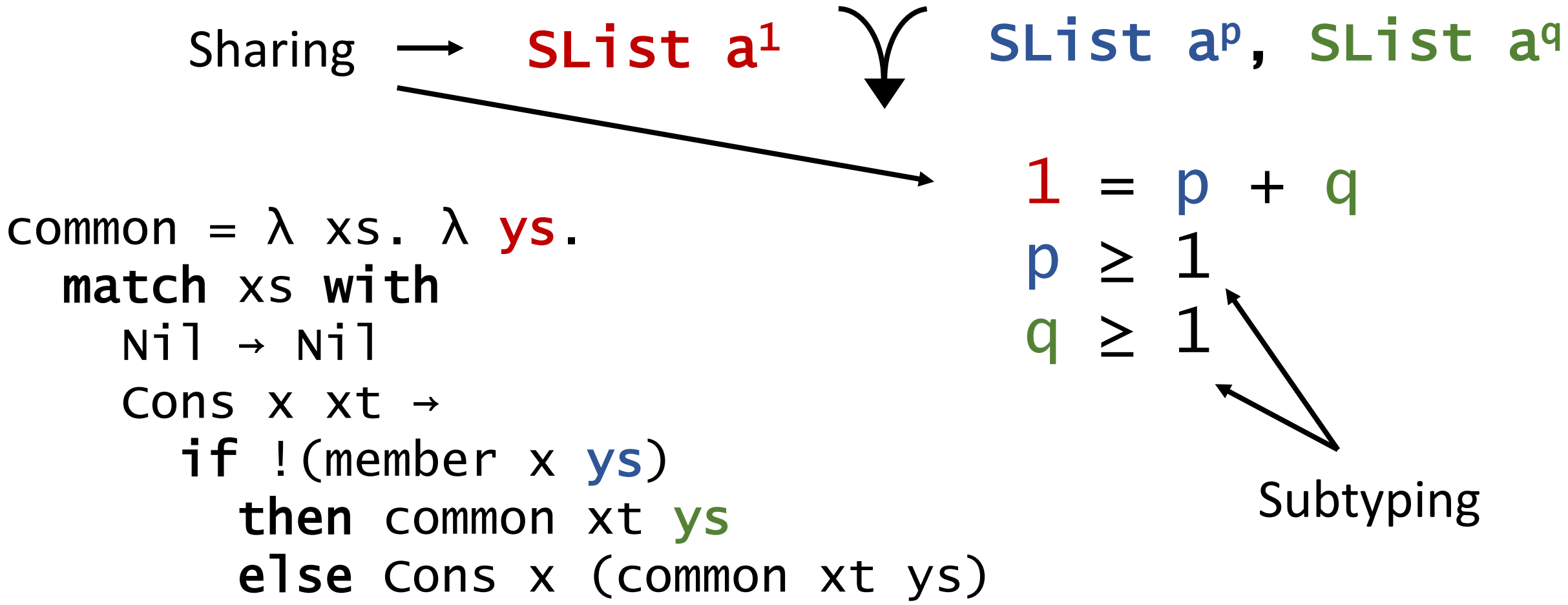


Sharing
SList a¹
SList aᵖ, SList aᵍ
1 = p + q
p ≥ 1
q ≥ 1
Subtyping
common = λ xs. λ ys.
match xs with
Nil → Nil
Cons x xt →
if !(member x ys)
then common xt ys
else Cons x (common xt ys)

SMT

$$1 = p + q \;\&\&\; p \geq 1 \;\&\&\; q \geq 1$$

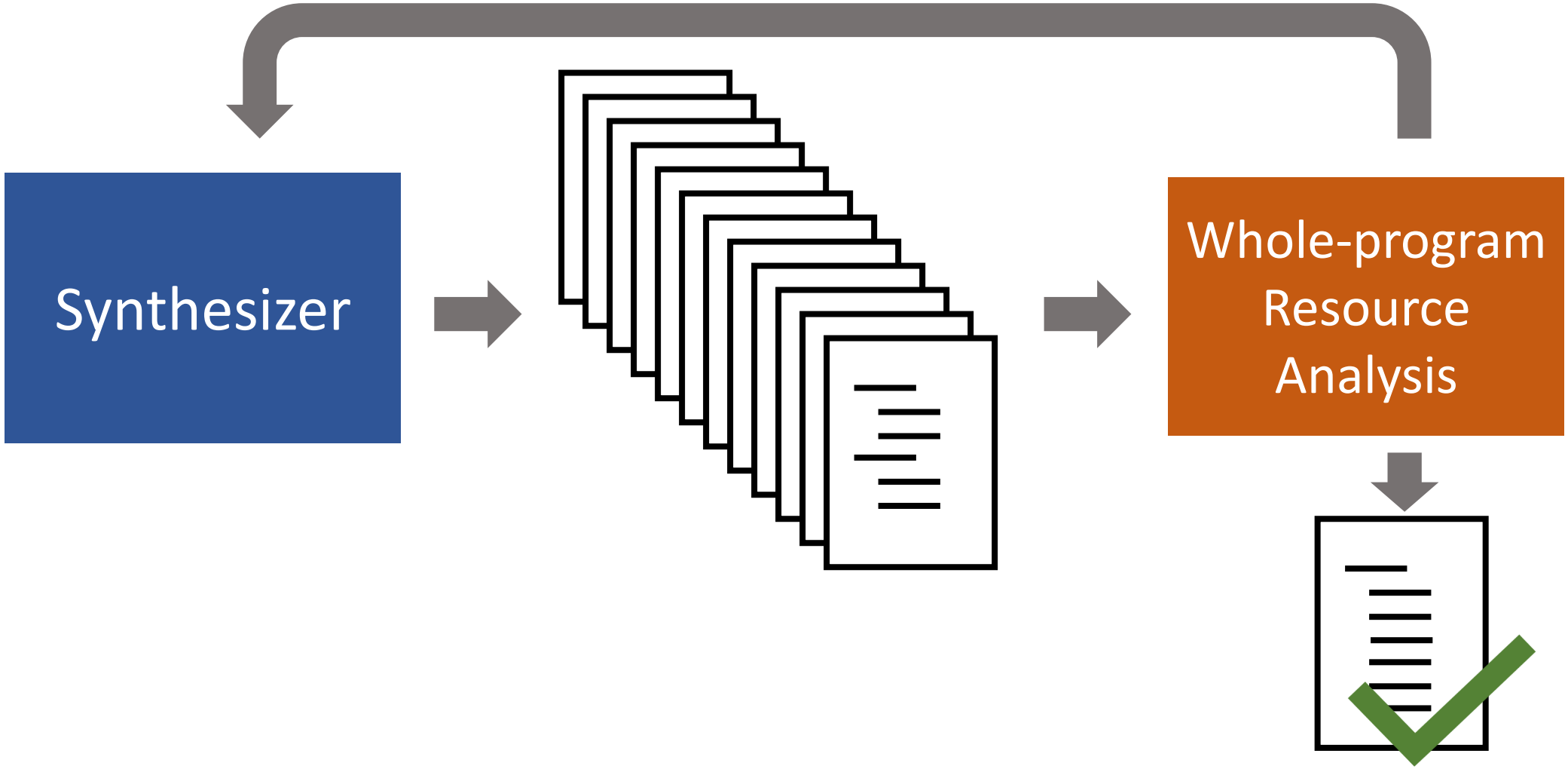
SMT

$$1 = p + q \;\&\&\; p \geq 1 \;\&\&\; q \geq 1$$

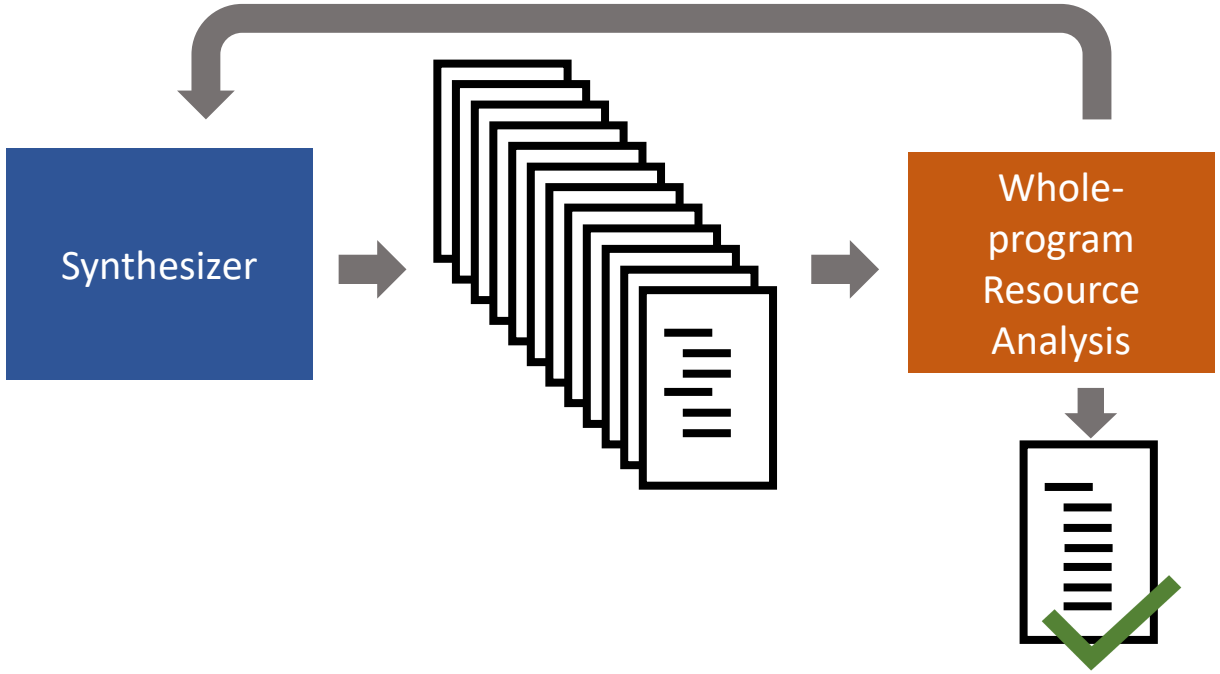
This talk

1. Specification
2. Analysis
3. Search

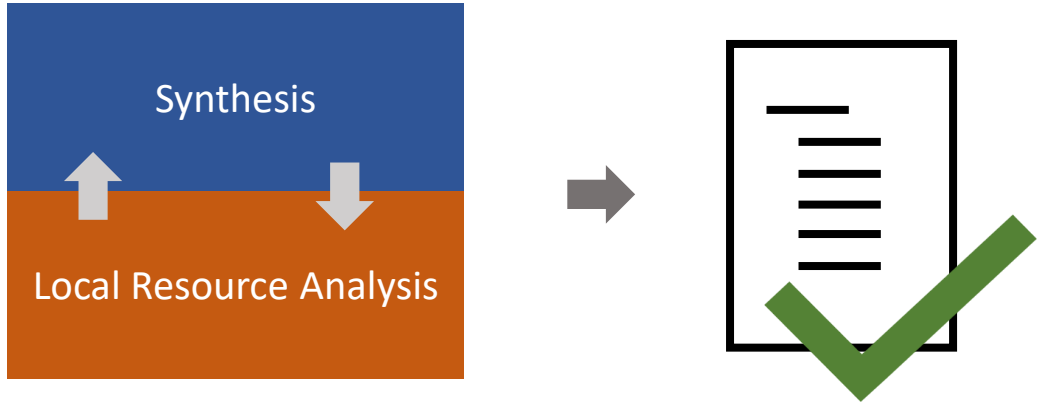
Enumerate-and-check



Enumerate-and-check



Resource-Guided Synthesis



Reject impossible programs early

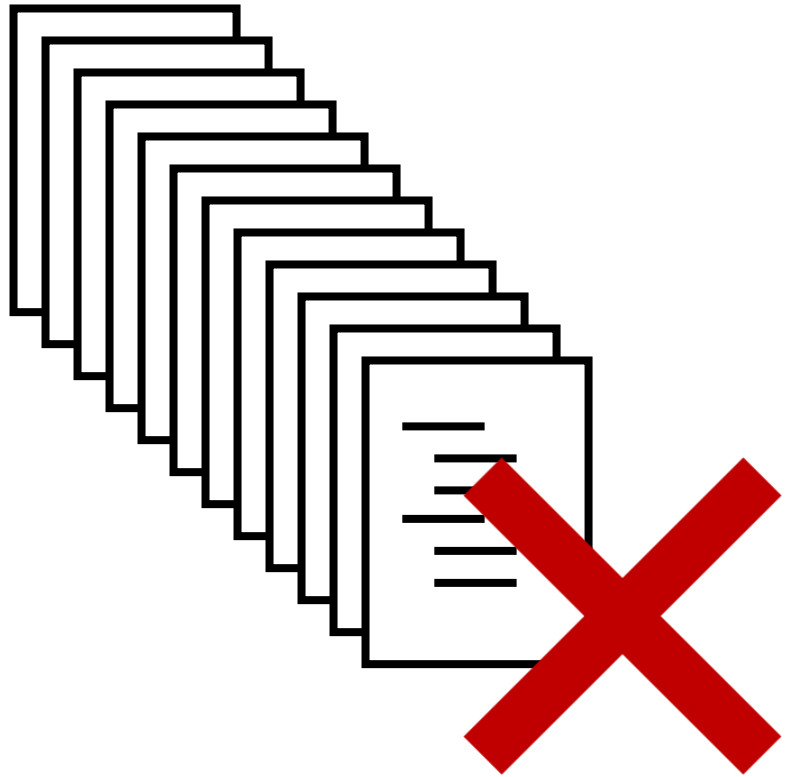
```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(member x ys)
        then common xt ys
        else ??
```

Reject impossible programs early with local analysis

```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(member x ys)
        then common xt ys
        else ??
```

Reject impossible programs early with local analysis

```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(mem x ys)
        then co  xt ys
        else
```



Reject impossible programs early with local analysis

```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(member x ys)
        then common ys ??
        else ??
```

Evaluation

Evaluation

1. Can ReSyn generate faster programs than Synquid?

Evaluation

1. Can ReSyn generate faster programs than Synquid?
2. How much longer does ReSyn take to generate code?

Evaluation

1. Can ReSyn generate faster programs than Synquid?
2. How much longer does ReSyn take to generate code?
3. Is local resource analysis effective at guiding the search?

1. Can ReSyn generate faster programs?

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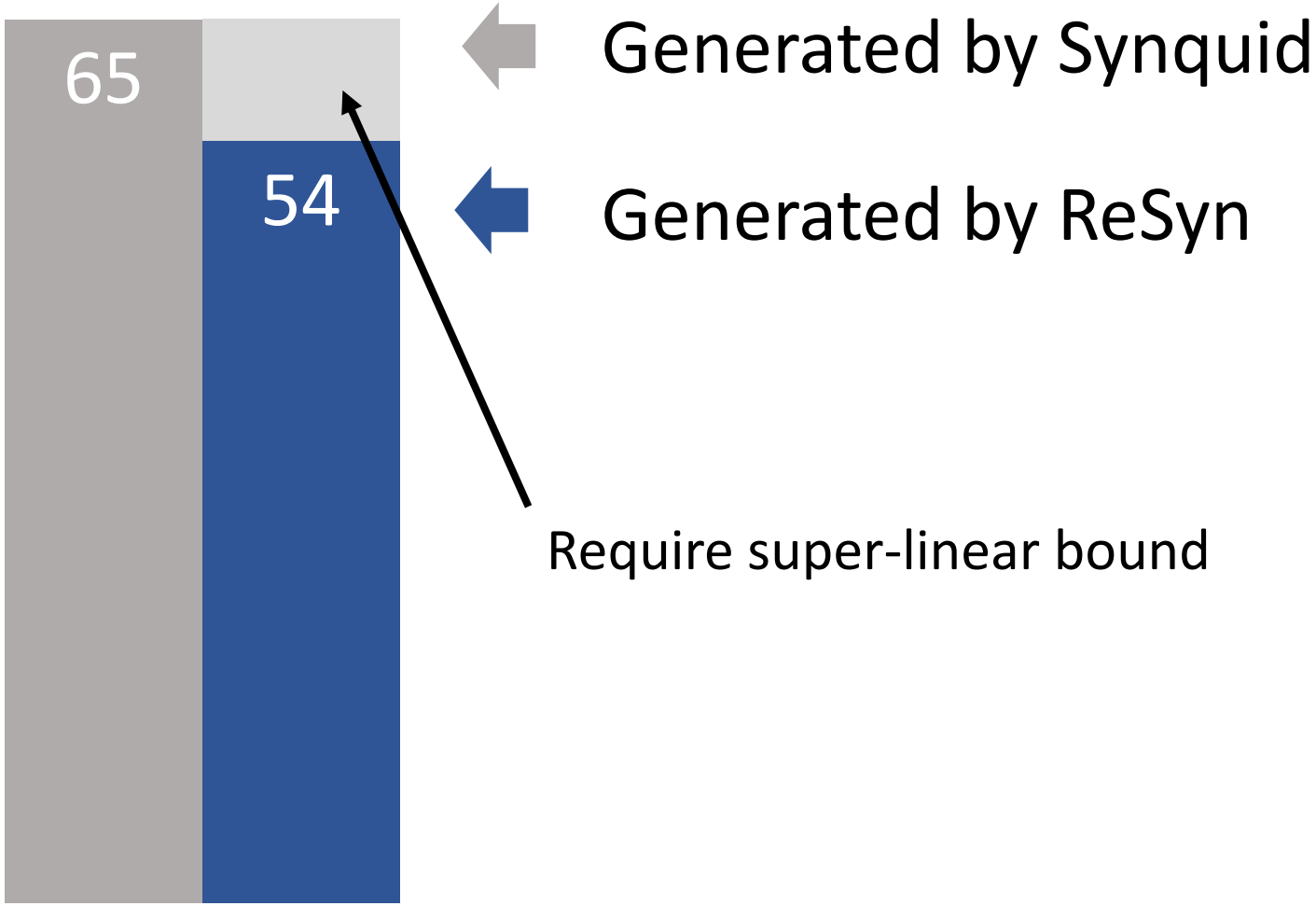
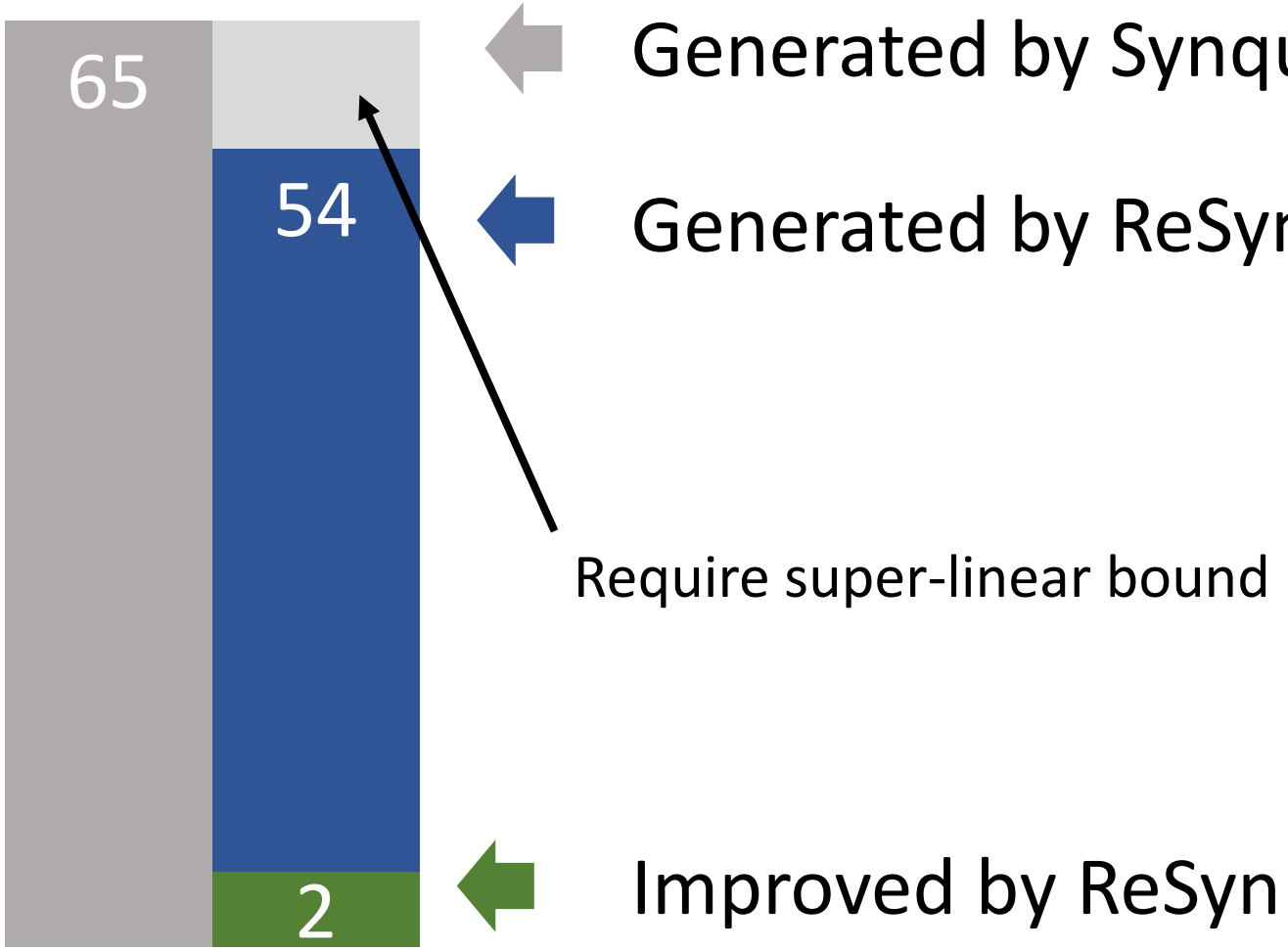


1. Can ReSyn generate faster programs?



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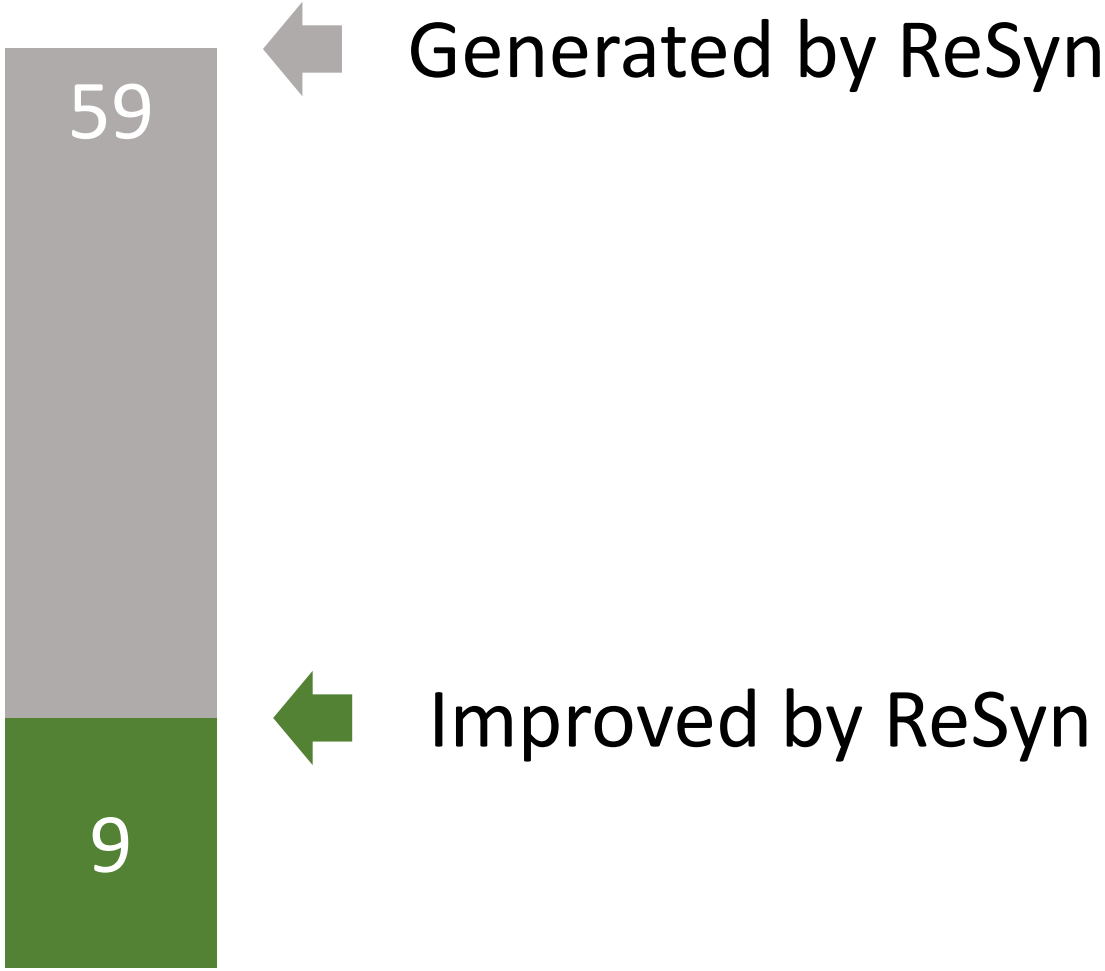
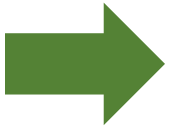


`compress`: Remove adjacent duplicates

```
compress xs =
  match xs with
    Nil → Nil
    Cons x3 x4 →
      match compress x4 with
        Nil → Cons x3 Nil
        Cons x10 x11 →
          if x3 == x10
            then compress x4
            else Cons x3 (Cons x10 x11)
```



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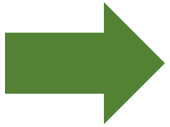
```

$O(2^n)$



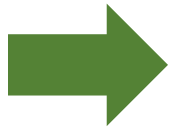
$O(n)$

Synquid

ReSyn

`insert`: Insert into a sorted list

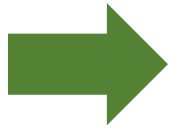
```
insert x xs =
  match xs with
    Nil → Cons x Nil
    Cons y ys →
      if x < y
        then Cons x (insert y ys)
        else Cons y (insert x ys)
```



```
insert x xs =
  match xs with
    Nil → Cons x Nil
    Cons y ys →
      if x < y
        then Cons x (Cons y ys)
        else Cons y (insert x ys)
```

O(n)

Synquid

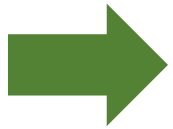


O(n)

ReSyn

```
insert :: x:a → xs: SList a^(if x > v then 1 else 0)
        → v:{SList a | elems v = elems xs ∪ {x}}
```

```
insert x xs =
  match xs with
    Nil → Cons x Nil
    Cons y ys →
      if x < y
        then Cons x (insert y ys)
        else Cons y (insert x ys)
```



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```
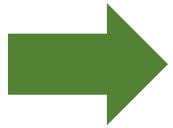
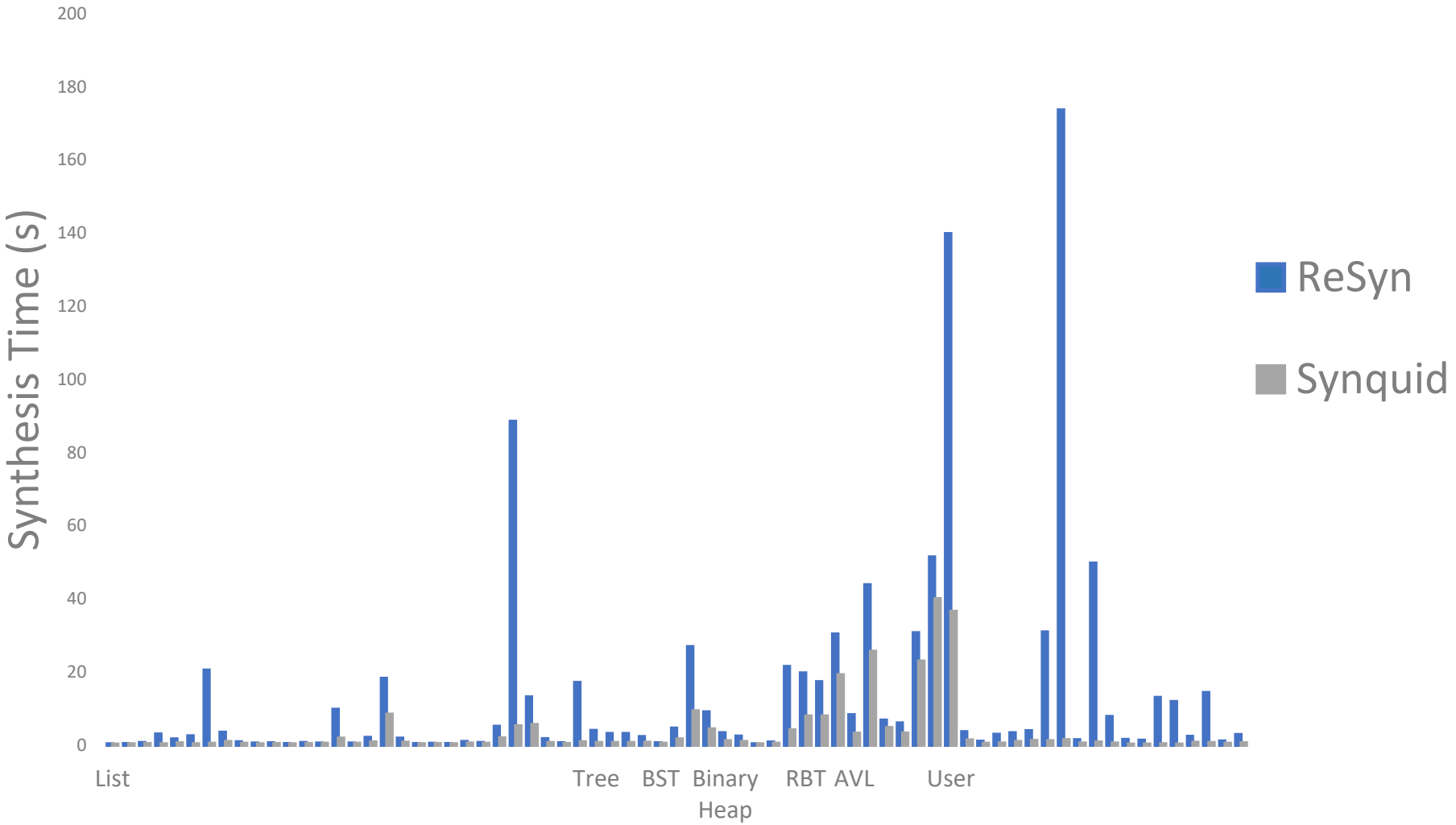

O(n)



O(n)

"One recursive call per element in xs that is smaller than x"

2. How do synthesis times compare?



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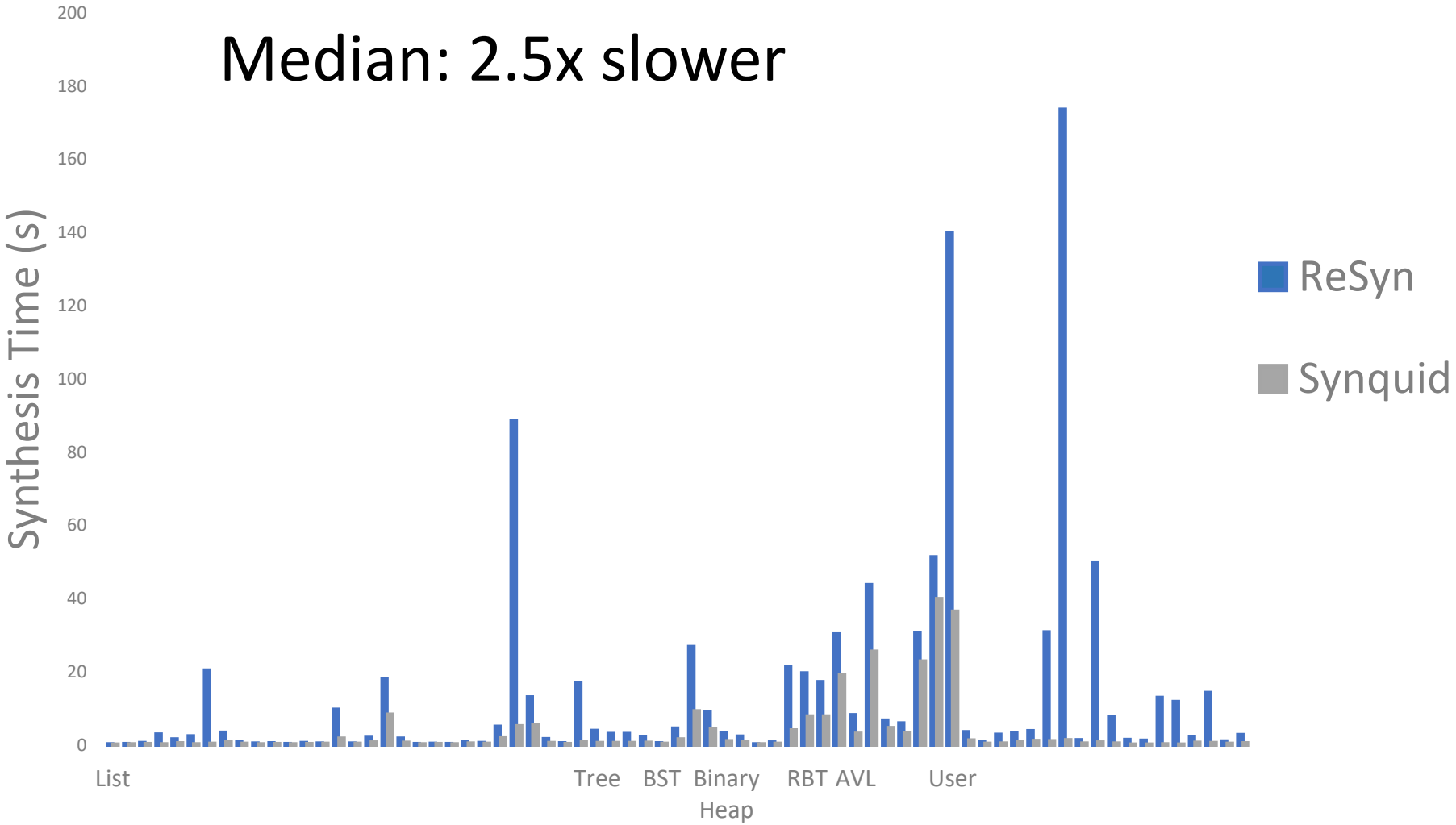

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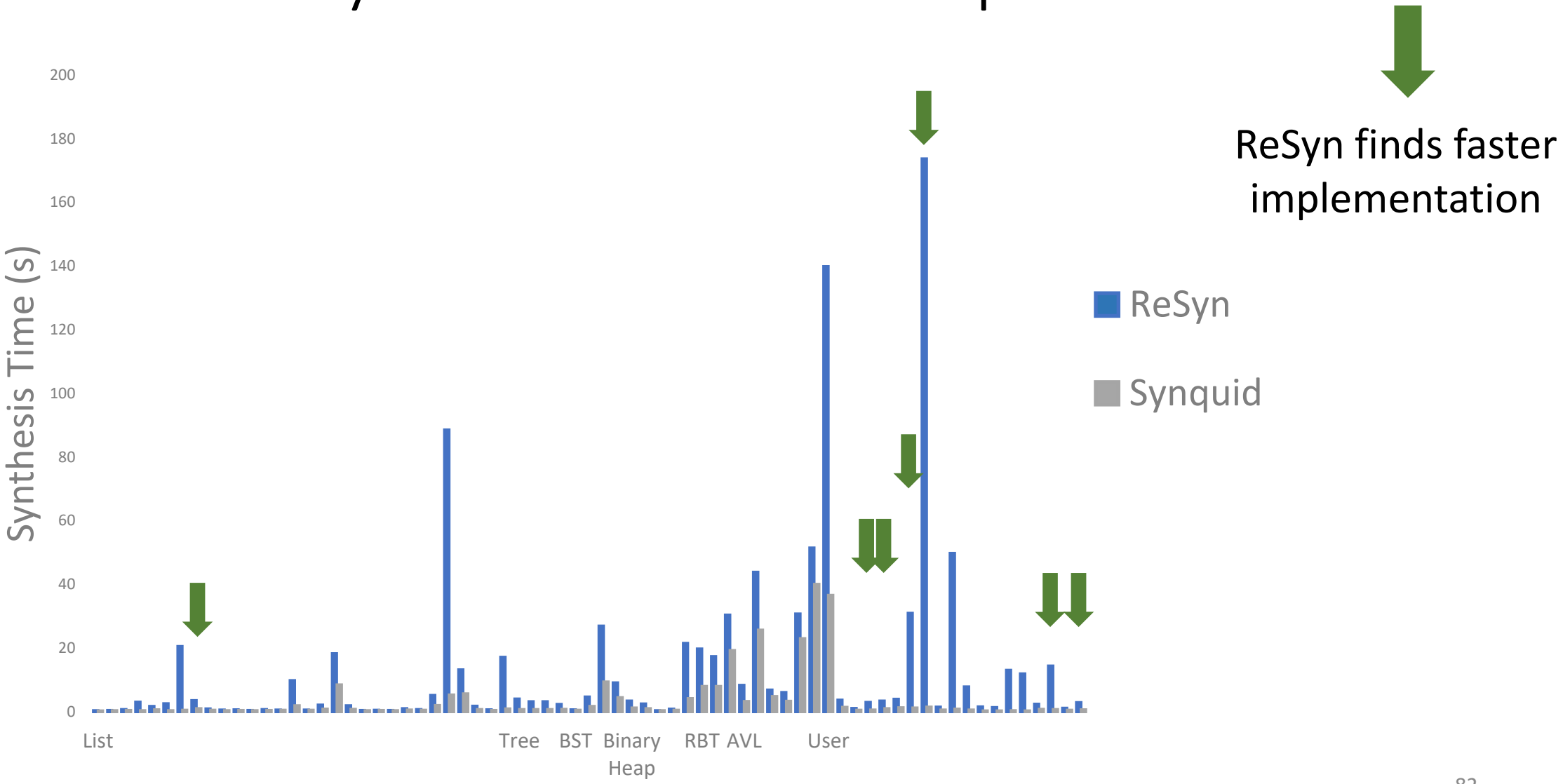

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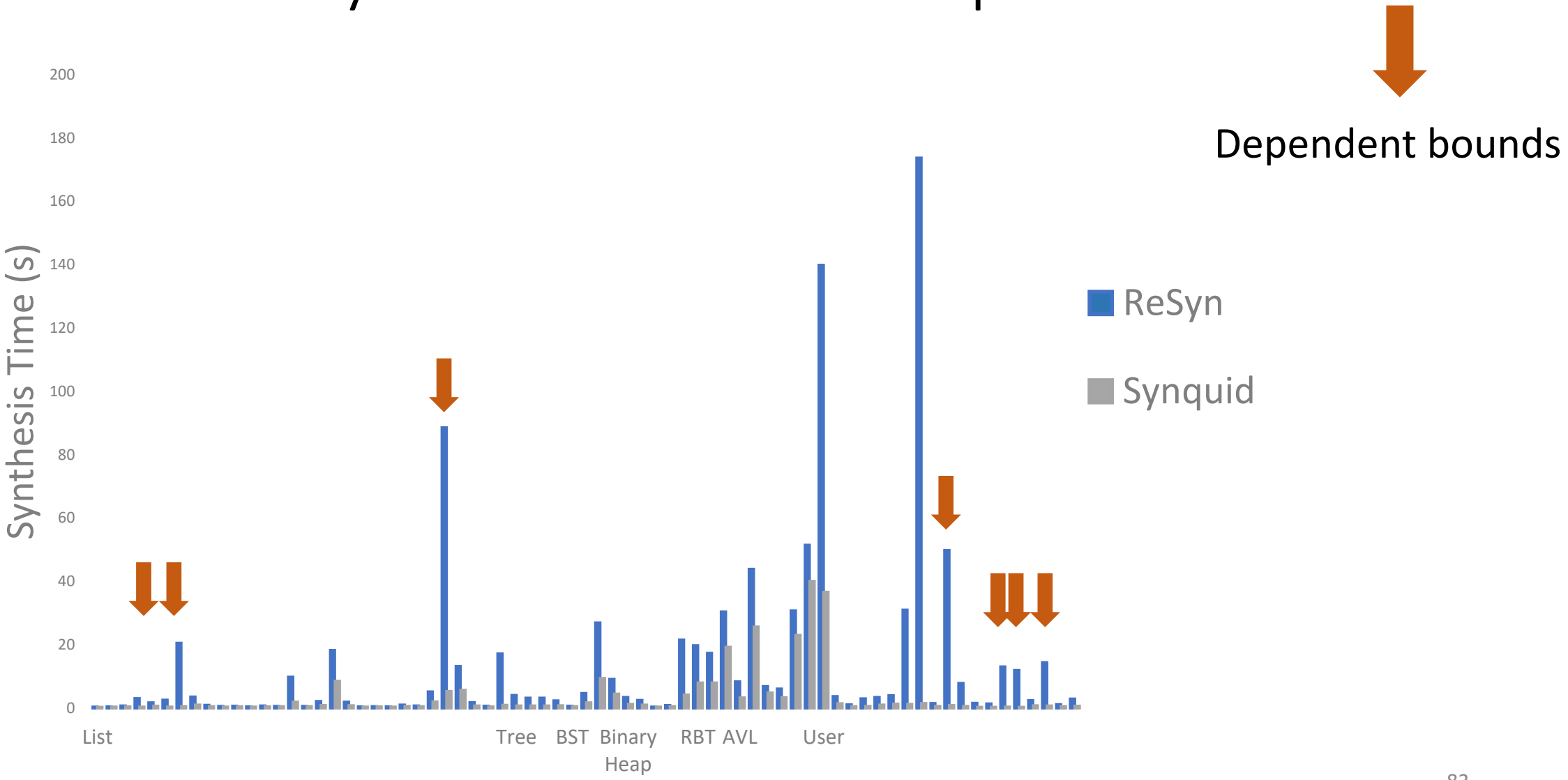

3. Does local resource analysis guide synthesis?

3. What happens if the analysis is non-local?

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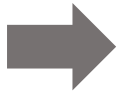
ReSyn

Enumerate-and-check

Timeout > 600s

3. What happens if the analysis is non-local?


```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      if !(member x ys)
        then common xt ys
        else Cons x (common x ys)
```

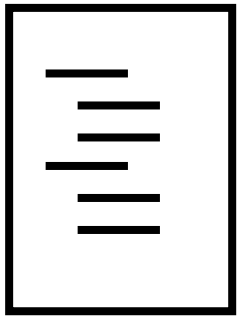


```
common = λ xs. λ ys.
  match xs with
    Nil → Nil
    Cons x xt →
      match ys with
        Nil → Nil
        Cons y yt →
          if x < y
            then common xt ys
            else if y < x
              then common xs yt
              else Cons x (common xs ys)
```

Timeout > 600s

What we had

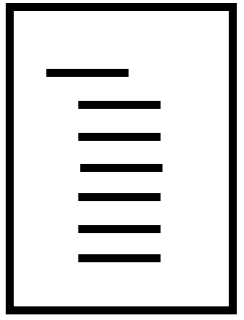
"Find the intersection of two sorted lists"



O(m·n)

What we have now

"Find the intersection of two sorted lists **in linear time**"



O(m+n)

https://bitbucket.org/tjknoth/resyn